

# Abstract Perturbed Krylov Methods

Just another point of view?

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# Outline

- 1 Getting started
  - the name of the game
  - a few examples
  - basic notations
  - HESSENBERG structure

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  - linear systems: (Q)OR
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We aim at 1a (possibly 3 and 4a), not 2.

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### Merriam-Webster Online: abstract (adjective)

- 1 a : disassociated from any specific instance
- 2 expressing a quality apart from an object
- 3 a : dealing with a subject in its abstract aspects

# perturbed KRYLOV methods

We consider perturbed KRYLOV subspace methods that can be written in the form

$$AQ_k = Q_{k+1} \underline{C}_k - F_k, \quad (1a)$$

$$Q_{k+1} \underline{C}_k = Q_k C_k + M_k, \quad (1b)$$

$$M_k = q_{k+1} c_{k+1,k} e_k^T. \quad (1c)$$

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We refer to the set of equations (1) as a *perturbed KRYLOV decomposition*.



# the main actors

In the perturbed KRYLOV decomposition:

- $A \in \mathbb{C}^{n \times n}$  is the system matrix from

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- $\underline{C}_k \in \mathbb{C}^{(k+1) \times k}$  is extended upper HESSENBERG
- $F_k \in \mathbb{C}^{n \times k}$  is zero or captures perturbations (due to finite precision, inexact methods, both, ...)

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We treat the system matrix  $A$ , the starting vector  $q_1$  and the perturbation terms  $\{f_l\}_{l=1}^k$  as input data and express everything else based on the *computed*  $C_k$ .

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# ARNOLDI

In the ARNOLDI method:

- $A \in \mathbb{C}^{n \times n}$  is a general matrix
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In the finite precision ARNOLDI method:

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- $C_k \in \mathbb{C}^{k \times k}$  is unreduced HESSENBERG
- $F_k \in \mathbb{C}^{n \times k}$  is “small”

(ask Miro about the details :-)

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In the inexact ARNOLDI method:

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In the LANCZOS method:

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The error terms may grow unbounded ...

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# power method

In the power method:

- $A \in \mathbb{C}^{n \times n}$  is a general matrix
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- $F_k \in \mathbb{C}^{n \times k}$  is “small” compared to  $Q_k$

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# a rather silly method

Consider any  $v \neq 0$  such that  $Av = v\lambda$  with  $\lambda \neq 0$

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Set

$$C_k \equiv \begin{pmatrix} o_{k-1}^T & 0 \\ \lambda I_{k-1} & \lambda e_{k-1} \end{pmatrix} \quad (2)$$

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Then  $AQ_k = Q_k C_k$ .

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# eigenmatrices et al.

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JORDAN matrices (, boxes) and blocks:

$$J_\Lambda = \oplus J_\lambda, \quad J_\lambda = \oplus J_{\lambda_\nu}, \quad J_\Theta = \oplus J_\theta. \quad (5)$$

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partial eigenmatrices:

$$V = \oplus V_\lambda, \quad V_\lambda = \oplus V_{\lambda\iota}, \quad S_k = \oplus S_\theta. \quad (6)$$



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reduced characteristic polynomial:

$$\chi_{C_k}(z) = (z - \theta)^{\alpha} \omega(z). \quad (10)$$

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# HESSENBERG eigenvalue-eigenmatrix relations

## Definition (off-diagonal products)

We denote the products of off-diagonal elements by

$$c_{i:j} \equiv \prod_{\ell=i}^j c_{\ell+1,\ell}. \quad (11)$$

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## Definition (polynomial vectors $\nu$ and $\check{\nu}$ )

We define vectors of (scaled) characteristic polynomials by

$$\nu(\mathbf{z}) \equiv \left( \frac{\chi_{C_{l+1:k}}(\mathbf{z})}{c_{l:k-1}} \right)_{l=1}^k, \quad \check{\nu}(\mathbf{z}) \equiv \left( \frac{\chi_{C_{l-1}}(\mathbf{z})}{c_{1:l-1}} \right)_{l=1}^k. \quad (12)$$

# HESSENBERG eigenvalue-eigenmatrix relations

## Definition (matrices of derivatives)

We define rectangular matrices collecting the derivatives by

$$\mathcal{S}_{\alpha-1}(\theta) \equiv \left[ \nu(\theta), \nu'(\theta), \frac{\nu''(\theta)}{2}, \dots, \frac{\nu^{(\alpha-1)}(\theta)}{(\alpha-1)!} \right] \quad (13)$$

$$\check{\mathcal{S}}_{\alpha-1}(\theta) \equiv \left[ \frac{\check{\nu}^{(\alpha-1)}(\theta)}{(\alpha-1)!}, \dots, \frac{\check{\nu}''(\theta)}{2}, \check{\nu}'(\theta), \check{\nu}(\theta) \right] \quad (14)$$



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## Observation

These matrices gather complete left and right JORDAN chains.

# HESSENBERG eigenvalue-eigenmatrix relations

## Theorem (HEER)

HESSENBERG *eigenmatrices* satisfy

$$\frac{P^{(\alpha-1)}(\theta)}{(\alpha-1)!} = S_\theta \omega(J_\theta) \check{S}_\theta^T = c_{1:k-1} S_{\alpha-1}(\theta) \check{S}_{\alpha-1}(\theta)^T. \quad (15)$$

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## Proof.

Proof based on comparison of TAYLOR expansions of the adjugate  $P(z)$  as inverse divided by determinant and the polynomial expression for the adjugate in terms of characteristic polynomials of submatrices (Zemke 2004, submitted to LAA). □

# HESSENBERG eigenvalue-eigenmatrix relations

## Lemma (HEER)

*We can choose the partial eigenmatrices such that*

$$\mathbf{e}_1^T \check{\mathbf{S}}_\theta = \mathbf{e}_\alpha^T (\omega(\mathbf{J}_\theta))^{-T}, \quad (16a)$$

$$\mathbf{S}_\theta^T \mathbf{e}_l = \mathbf{c}_{1:l-1} \chi_{\mathbf{C}_{l+1:k}}(\mathbf{J}_\theta)^T \mathbf{e}_1. \quad (16b)$$

*Tailored to diagonalizable  $\mathbf{C}_k$ :*

$$\check{\mathbf{S}}_{lj} \mathbf{S}_{\ell j} = \frac{\chi_{\mathbf{C}_{1:l-1}}(\theta_j) \mathbf{c}_{l:\ell-1} \chi_{\mathbf{C}_{\ell+1:k}}(\theta_j)}{\chi'_{\mathbf{C}_k}(\theta_j)} \quad \forall l \leq \ell. \quad (17)$$

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# basic definitions

## Definition (basis polynomials)

We define the (trailing) basis polynomials by

$$B_k(\mathbf{z}) \equiv \frac{\chi_{C_k}(\mathbf{z})}{C_{1:k}} = \check{\nu}_{k+1}(\mathbf{z}), \quad (18)$$

$$B_{l+1:k}(\mathbf{z}) \equiv \frac{\chi_{C_{l+1:k}}(\mathbf{z})}{C_{l+1:k}} = \frac{C_{l+1,l}}{C_{k+1,k}} \nu_l(\mathbf{z}), \quad \forall l = 1, \dots, k. \quad (19)$$

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## Observation

The trailing basis polynomials are the basis polynomials of the trailing submatrices  $C_{l+1:k}$ .

# "basis" vectors

## Theorem (the "basis" vectors)

*The "basis" vectors of a KRYLOV method are given by*

$$q_{k+1} = \mathcal{B}_k(A)q_1 \quad . \quad (20)$$



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## Observation

The **perturbed** "basis" vectors can be interpreted as an additive overlay of **exact** "basis" vectors.

# a rough sketch of a short proof

## Proof.

Introduce variable  $z$ :

$$M_k = Q_k(zI - C_k) + (zI - A)Q_k + F_k$$

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HEER:  $\operatorname{adj}(zC_k)e_1 = c_{1:k-1}\nu(z)$ . Insert  $A$  into

$$c_{k+1,k}q_{k+1} = \frac{q_1 \chi_{C_k}(z)}{c_{1:k-1}} + (zI - A)Q_k\nu(z) + F_k\nu(z). \quad \square$$

# a closer & deeper look

## Theorem (the "basis" vectors revisited)

Let  $C_k$  be diagonalizable and suppose that  $\lambda \neq \theta_j$  for all  $j$ :

$$\left( \sum_{j=1}^k \frac{c_{1:k}}{\chi'_{C_k}(\theta_j)(\lambda - \theta_j)} \right) \hat{v}^H q_{k+1} = \hat{v}^H q_1$$

# a closer & deeper look

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## Remark

Generalization to the non-diagonalizable case exists.

# Outline

- 1 Getting started
  - the name of the game
  - a few examples
  - basic notations
  - HESSENBERG structure
- 2 **The results ...**
  - "basis" transformations
  - **eigenvalue problems**
  - linear systems: (Q)OR
  - linear systems: (Q)MR
- 3 ... and their impacts
  - general comments
  - finite precision issues
  - inexact KRYLOV methods

# eigenvalues, JORDAN block, partial eigenmatrix

Unreduced HESSENBERG matrices  $C_k$  are non-derogatory.

## Notations

In the following,

The matrices are such that

$$C_k S_\theta = S_\theta J_\theta, \quad \text{where } J_\theta \in \mathbb{C}^{\alpha \times \alpha}. \quad (21)$$

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# RITZ pairs, RITZ residuals

## Definition (RITZ pair)

Define RITZ pair by

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## Observation

A backward expression for the RITZ residual is given by

$$AY_\theta - Y_\theta J_\theta = q_{k+1} c_{k+1,k} e_k^T S_\theta - F_k S_\theta. \quad (23)$$

# RITZ residuals (generic case)

## Theorem (generic RITZ residuals)

*The RITZ residual for an (arbitrarily chosen) RITZ pair:*

$$\begin{aligned}
 AY_\theta - Y_\theta J_\theta = & \left( \frac{\chi_{C_k}(A)}{c_{1:k}} \right) q_1 e_k^T S_\theta \\
 & + \sum_{l=1}^k \left( \frac{\chi_{C_{l+1:k}}(A)}{c_{l:k-1}} \right) f_l e_k^T S_\theta - f_l e_l^T S_\theta. \quad (24)
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## Proof.

Backward expression and Theorem on the "basis" vectors. □

## RITZ residuals (special case)

Use (unique) choice for the partial eigenmatrix  $S_\theta$  (HEER):

### Theorem (special RITZ residuals)

*The RITZ residual for the special partial eigenmatrix from HEER is given by*

$$\begin{aligned}
 AY_\theta - Y_\theta J_\theta &= \chi_{C_k}(A)q_1 e_1^T \\
 &+ \sum_{l=1}^k c_{1:l-1} \left( \chi_{C_{l+1:k}}(A)f_l e_1^T - f_l e_1^T \chi_{C_{l+1:k}}(J_\theta) \right). \quad (25)
 \end{aligned}$$

# bivariate adjugate polynomials

## Definition (bivariate adjugate polynomials)

We define the bivariate adjugate polynomials by

$$\mathcal{A}_k(\theta, z) \equiv \begin{cases} (\chi_{C_k}(\theta) - \chi_{C_k}(z)) (\theta - z)^{-1}, & z \neq \theta, \\ \chi'_{C_k}(z), & z = \theta. \end{cases} \quad (26)$$

Trailing bivariate adjugate polynomials  $\mathcal{A}_{l+1:k}$  are defined using  $C_{l+1:k}$  in place of  $C_k$ ,  $l = 1, \dots, k$ .

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## Observation

Even with an eigenvalue  $\theta$ :  $\mathcal{A}_k(\theta, C_k) = \text{adj}(\theta I_k - C_k) = P(\theta)$ .

# RITZ vectors

## Theorem (the RITZ vectors)

The RITZ vectors of a KRYLOV method are given by

$$\text{vec}(Y_\theta) = \begin{pmatrix} \mathcal{A}_k(\theta, A) \\ \mathcal{A}'_k(\theta, A) \\ \vdots \\ \frac{\mathcal{A}_k^{(\alpha-1)}(\theta, A)}{(\alpha-1)!} \end{pmatrix} q_1 \quad . \quad (27)$$

(derivation with respect to "shift"  $\theta$ )



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The RITZ vectors of a perturbed KRYLOV method are given by

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# sketch of proof: basics

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- The adjugate of a matrix is defined as matrix of cofactors.
- The adjugate is linked to eigenvectors and, more general, principal vectors.
- The adjugate is linked to the inverse and the determinant.

The problem: the definition of the bivariate adjugate polynomials given here is not "adequate", we need another form.

# sketch of proof: HESSENBERG basics

To derive this peculiar form we use the first adjugate identity:

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*First adjugate identity:*

$$(z - \theta)\text{adj}({}^zA)\text{adj}({}^\theta A) = \det({}^zA)\text{adj}({}^\theta A) - \det({}^\theta A)\text{adj}({}^zA). \quad (28)$$



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*Specialized to HESSENBERG matrices:*

$$(z - \theta) \sum_{j=1}^k \chi_{C_{1:j-1}}(z) \chi_{C_{j+1:k}}(\theta) = \chi_{C_k}(z) - \chi_{C_k}(\theta). \quad (29)$$

# sketch of proof: gluing results together

The last line implies the following representations ( $\ell \geq 0$ ):

$$\mathcal{A}_{l+1:k}^{(\ell)}(\theta, \mathbf{z}) = \sum_{j=l+1}^k \chi_{\mathbf{C}_{l+1:j-1}}(\mathbf{z}) \chi_{\mathbf{C}_{j+1:k}}^{(\ell)}(\theta) \quad \forall l = 0, 1, \dots, k. \quad (30)$$

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This together with

- the special choice of the partial eigenmatrix  $S_\theta$
- the representation of the "basis" vectors

are the building blocks for the proof.

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- 1 Getting started
  - the name of the game
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## (Q)OR: the approach

Suppose that  $C_k$  is invertible and that  $q_1 = r_0/\|r_0\|$ . Let  $z_k$  denote the solution to the linear system of equations

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$$x_k = Q_k z_k \quad (32)$$

and the  $k$ th (true) (Q)OR residual by

$$r_k = r_0 - Ax_k. \quad (33)$$

# a backward expression for the (Q)OR residual

## Observation

A backward expression for the (Q)OR residual is given by

$$r_k = r_0 - Ax_k = (Q_k C_k - A Q_k) C_k^{-1} e_1 \|r_0\|$$

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$$\begin{aligned}
 r_k &= r_0 - Ax_k = (Q_k C_k - A Q_k) C_k^{-1} e_1 \|r_0\| \\
 &= (-q_{k+1} c_{k+1,k} e_k^T + F_k) z_k
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 &= -q_{k+1} c_{k+1,k} z_{kk} + \sum_{l=1}^k f_l z_{lk}
 \end{aligned}$$

# adjugate, inverse, determinant

Express the inverse of  $C_k$  as adjugate by determinant:

$$\frac{-z_{lk}}{\|r_0\|} = e_l^T (-C_k)^{-1} e_1$$

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Utilize

$$r_k = q_{k+1} c_{k+1,k}(-z_{kk}) - \sum_{l=1}^k f_l(-z_{lk}). \quad (34)$$



## (Q)OR: the residuals

This backward expression plus Theorem on the "basis" vectors:

### Theorem (the (Q)OR residual vectors)

*The residual vectors of a*

*(Q)OR KRYLOV method are*

*given by*

$$r_k = \frac{\chi_{C_k}(A)}{\chi_{C_k}(0)} r_0 \quad . \quad (35)$$

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This backward expression plus Theorem on the "basis" vectors:

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The perturbation terms remind of adjugate polynomials ...

# adjugate, inverse, interpolation (I)

## Definition (univariate adjugate polynomials)

We define univariate adjugate polynomials by

$$\mathcal{A}_k(z) = (-1)^k (\chi_{C_k}(0) - \chi_{C_k}(z)) z^{-1}$$

By CAYLEY-HAMILTON:  $\mathcal{A}_k(C_k) = \text{adj}(C_k)$

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## Observation

Univariate and bivariate adjugate polynomials are related by

$$\mathcal{A}_k(z) = (-1)^{k+1} \mathcal{A}_k(z, 0) = (-1)^{k+1} \mathcal{A}_k(0, z)$$

# adjugate, inverse, interpolation (II)

## Notations

We define and denote the LAGRANGE interpolation of the inverse by

$$\mathcal{L}_k[z^{-1}](z) = \frac{A_k(z)}{\det(C_k)} = \left(1 - \frac{\chi_{C_k}(z)}{\chi_{C_k}(0)}\right) z^{-1}$$

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We define and denote the LAGRANGE interpolation of a perturbed identity by

$$\mathcal{L}_k^0[1 - \delta_{z0}](z) = \mathcal{L}_k[z^{-1}](z)z = \frac{\chi_{C_k}(0) - \chi_{C_k}(z)}{\chi_{C_k}(0)}.$$

# trailing {adjugate, inverse, interpolation}

We expand all notations to the trailing submatrices  $C_{l+1:k}$ .



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# (Q)OR: the residuals

## Theorem (the (Q)OR residual vectors)

*Suppose that all submatrices  $C_{l+1:k}$  are nonsingular.  
Then the residual vectors can be written as*

$$r_k = \frac{\chi_{C_k}(A)}{\chi_{C_k}(0)} r_0 - \sum_{l=1}^k z_{lk} \mathcal{L}_{l+1:k}^0 [1 - \delta_{z0}](A) f_l. \quad (37)$$

This occurs frequently, consider e.g. CG for HPD  $A$ .

## (Q)OR: the errors, regular $A$

What about the error vectors?



# (Q)OR: the errors, regular $A$

What about the error vectors?

**Theorem (the (Q)OR error vectors, regular  $A$ )**

*Suppose that  $A$  is invertible and let  $x = A^{-1}r_0$  denote the unique solution of the linear system  $Ax = r_0$ .*

*Then the error vectors are given by*

$$(x - x_k) = \frac{\chi_{C_k}(A)}{\chi_{C_k}(0)}(x - 0) + \|r_0\| \sum_{l=1}^k c_{1:l-1} \frac{A_{l+1:k}(0, A)}{\chi_{C_k}(0)} f_l. \quad (38)$$

## (Q)OR: the errors, regular $A$

What about invertible submatrices?

# (Q)OR: the errors, regular $A$

What about invertible submatrices?

Theorem (the (Q)OR error vectors, regular  $A$  and  $C_{l+1:k}$ )

*Suppose that all trailing submatrices  $C_{l+1:k}$  are nonsingular. Then the error vectors can be written as*

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# (Q)OR: the errors, singular $A$

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Theorem (the (Q)OR error vectors, singular  $A$ )

When  $A$  is singular, with  $x \equiv A^D r_0$ , where  $A^D$  denotes the DRAZIN inverse of  $A$ ,

$$(x - AA^D x_k) = \frac{\chi_{C_k}(A)}{\chi_{C_k}(0)} (x - 0) + \|r_0\| \sum_{l=1}^k c_{1:l-1} \frac{A_{l+1:k}(0, A)}{\chi_{C_k}(0)} AA^D f_l. \quad (40)$$

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Theorem (the (Q)OR error vectors, singular A, regular  $C_{l+1:k}$ )

When A is singular, with  $x \equiv A^D r_0$ ,

$$(x - AA^D x_k) = \frac{\chi_{C_k}(A)}{\chi_{C_k}(0)} (x - 0) - \sum_{l=1}^k z_{lk} \mathcal{L}_{l+1:k}[z^{-1}](A) AA^D f_l. \quad (41)$$

## (Q)OR: the iterates

The iterates  $x_k$  can be composed like the RITZ vectors.



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### Theorem (the (Q)OR iterates)

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Theorem (the (Q)OR iterates, regular  $C_{l+1:k}$ )

*Suppose that all  $C_{l+1:k}$  are regular.*

*Then*

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### Observation

This is a linear combination of  $k + 1$  approximations from distinct KRYLOV subspaces, spanned by the same matrix  $A$ , *but distinct starting vectors.*

# Outline

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  - the name of the game
  - a few examples
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  - **linear systems: (Q)MR**
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## (Q)MR: the approach

Let  $\underline{z}_k$  denote the minimal-norm solution of the least-squares problem

$$\|\underline{C}_k \underline{z}_k - \underline{e}_1\|_{r_0} = \min. \quad (44)$$

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Define the  $k$ th (Q)MR iterate  $\underline{x}_k$  by

$$\underline{x}_k = Q_k \underline{z}_k \quad (45)$$

and the  $k$ th quasi-residual by

$$\tau_k = \underline{e}_1 \|r_0\| - \underline{C}_k \underline{z}_k = (\underline{I}_k - \underline{C}_k \underline{C}_k^\dagger) \underline{e}_1 \|r_0\|. \quad (46)$$



# (Q)MR: a backward expression for the residual

## Observation

The residual  $\underline{r}_k$  of the (Q)MR iterates has the following backward expression:

$$\underline{r}_k = r_0 - A\underline{x}_k = Q_{k+1}\underline{e}_1\|r_0\| - AQ_k\underline{z}_k \quad (47)$$

$$= Q_{k+1}(\underline{e}_1\|r_0\| - \underline{C}_k\underline{z}_k) + F_k\underline{z}_k = Q_{k+1}\underline{\tau}_k + \sum_{l=1}^k f_l\underline{z}_{lk}. \quad (48)$$

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## Observation

To express the residual  $\underline{r}_k$  as polynomial in  $A$ , we "only" need "polynomial" expressions for  $\underline{\tau}_k$  and  $\underline{z}_k$ .

# (Q)MR: HESSENBERG rewritings

Definition (the scalar vectors  $\mu$ ,  $\check{\mu}$  and  $\hat{\mu}$ )

We define pairs of vectors  $\mu^j, \check{\mu}^j \in \mathbb{C}^j$  and  $\hat{\mu}^j \equiv \overline{\check{\mu}^j} \in \mathbb{C}^j$ :

$$\mu \equiv \left( \frac{(-1)^{l+1} \det(\mathbf{C}_{l+1:j})}{c_{l:j-1}} \right)_{l=1}^j, \quad (49)$$

$$\check{\mu} \equiv \left( \frac{(-1)^{j-l} \det(\mathbf{C}_{l-1})}{c_{1:l-1}} \right)_{l=1}^j. \quad (50)$$

# (Q)MR: HESSENBERG rewritings

## Lemma (MOORE-PENROSE inverse of extended HESSENBERG)

The MOORE-PENROSE inverse of the extended HESSENBERG matrix  $\underline{C}_k$  is given by

$$\underline{C}_k^\dagger = \frac{\sum_{j=1}^k |c_{j+1:k}|^2 \begin{pmatrix} \overline{\det(C_j)} \text{adj}(C_j) & \overline{c_{1:j}} \text{adj}(C_j) \hat{\mu}^j & O_{j,k-j} \\ O_{k-j,j} & o_{k-j} & O_{k-j} \end{pmatrix}}{\sum_{j=0}^k |c_{j+1:k}|^2 |\det(C_j)|^2}. \quad (51)$$

# (Q)MR: HESSENBERG rewritings

## Lemma (the minimal norm solution)

The minimal norm solution  $\underline{z}_k$  is given by

$$\frac{\underline{z}_k}{\|r_0\|} = \frac{\sum_{j=1}^k |c_{j+1:k}|^2 \begin{pmatrix} \overline{\det(C_j)} c_{1:j-1} \mu^j \\ o_{k-j} \end{pmatrix}}{\sum_{j=0}^k |c_{j+1:k}|^2 |\det(C_j)|^2} \quad (52)$$

$$= (-1)^{k+1} \frac{\begin{pmatrix} o_k & \text{adj}(C_{k+1}^\Delta) \end{pmatrix} \text{adj}(C_{k+1}^H) e_{k+1}}{\sum_{j=0}^k |c_{j+1:k}|^2 |\det(C_j)|^2}. \quad (53)$$

# (Q)MR: HESSENBERG rewritings

## Lemma ((Q)MR and (Q)OR)

Suppose all leading  $C_j$  are regular. Then the relation between the  $k$ th (Q)MR solution  $\underline{z}_k$  and all prior (Q)OR solutions  $z_j$  is given by

$$\underline{z}_k = \frac{\sum_{j=0}^k |\det(C_j)|^2 |c_{j+1:k}|^2 \begin{pmatrix} z_j \\ \mathbf{o}_{k-j} \end{pmatrix}}{\sum_{j=0}^k |\det(C_j)|^2 |c_{j+1:k}|^2}, \quad (54)$$

where  $z_0$  is the empty matrix with dimensions  $0 \times 1$ .

# (Q)MR: HESSENBERG rewritings

## Lemma (the quasi-residual)

The quasi-residual  $\tau_k$  is given by

$$\frac{\tau_k}{\|r_0\|} = c_{1:k} \left( \frac{(-1)^{l-1} c_{l:k} \det(C_{l-1})}{\sum_{j=0}^k |c_{j+1:k}|^2 |\det(C_j)|^2} \right)_{l=1}^{k+1}. \quad (55)$$

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Lacking is the "right" **interpretation**.

This is currently work in progress.

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- are **aiding** the design of particular finite precision/inexact methods.

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### Example (inexact (Q)OR, e.g., inexact CG)

We have proven

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Based on the behavior of the solution vectors  $z_k$  and/or the LAGRANGE interpolations we can allow the perturbation vectors  $f_l$  to grow (in certain directions).

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- (Q)MR case has to be investigated more thoroughly.

that's all ...

Děkuji.