# Abstract Perturbed Krylov Methods 

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## Outline

## Philosophical considerations

A matrix equation The iterative point of view The polynomial point of view

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The results on ...
Ritz vectors
QOR iterates
QOR residuals

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## A "numerical" experiment

Eigenvectors using Lanczos' method

## A matrix-theoretical beginning

We start with the equation

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\begin{aligned}
A Q_{k} \quad & =Q_{k+1} C_{k} \\
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Q_{k+1}=\left(\begin{array}{ll}
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F_{k} & \in \mathbb{C}^{(n, k)} & \text { is a general perturbation. }
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## The dependence on the iteration

We investigate this matrix equation iteratively:

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A Q_{k} e_{l}+F_{k} e_{l} & =Q_{k+1} C_{k} e_{l} \\
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Idea:
- Interpret perturbed Krylov methods as overlay of several polynomial methods.


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- residual polynomials $\mathcal{R}_{k}$ and $\underline{\mathcal{R}}_{k}$.

We restrict ourselves to $\mathcal{A}_{k}, \mathcal{L}_{k}\left[z^{-1}\right], \mathcal{L}_{k}\left[1-\delta_{z 0}\right]$ and $\mathcal{R}_{k}$.

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- This implies [Z, 2006]

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\mathcal{A}_{k}\left(\theta_{j}, C_{k}\right) e_{1}=s_{j}, \quad C_{k} s_{j}=\theta_{j} s_{j}
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- Generalization:

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\mathcal{A}_{l+1: k}(\theta, z) \equiv \frac{\chi_{l+1: k}(\theta)-\chi_{l+1: k}(z)}{\theta-z}, \quad l=0,1, \ldots, k
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## Adjugate polynomials and Ritz vectors

## Theorem (Ritz vectors)

Let $C_{k} S_{\theta}=S_{\theta} J_{\theta}$ (for a certain $S_{\theta}$ ). Let the Ritz matrix be given by $Y_{\theta} \equiv Q_{k} S_{\theta}$. Then

$$
\operatorname{vec}\left(Y_{\theta}\right)=\left(\begin{array}{c}
\mathcal{A}_{k}(\theta, A)  \tag{3}\\
\mathcal{A}_{k}^{\prime}(\theta, A) \\
\vdots \\
\frac{\mathcal{A}_{k}^{(\alpha-1)}(\theta, A)}{(\alpha-1)!}
\end{array}\right) q_{1}+\sum_{l=1}^{k} c_{1: l-1}\left(\begin{array}{c}
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\end{array}\right) f_{l,}
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We might scale differently such that (here only for approximate eigenvectors)

$$
y=\frac{\mathcal{A}_{k}(\theta, A)}{\prod_{\ell=1}^{k-1} c_{\ell+1, \ell}} q_{1}+\sum_{l=1}^{k} \frac{\mathcal{A}_{l+1: k}(\theta, A)}{\prod_{\ell=l+1}^{k-1} c_{\ell+1, \ell}} \frac{f_{l}}{c_{l+1, l}}
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## Lagrange polynomials and QOR iterates

## Theorem (QOR iterates)

Suppose all $C_{l+1: k}$ are regular. Define $z_{k} \equiv C_{k}^{-1} e_{1}\left\|r_{0}\right\|$ and $x_{k} \equiv Q_{k} z_{k}$. Then

$$
\begin{equation*}
x_{k}=\mathcal{L}_{k}\left[z^{-1}\right](A) r_{0}-\sum_{l=1}^{k} z_{l k} \mathcal{L}_{l+1: k}\left[z^{-1}\right](A) f_{l} \tag{4}
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Really sloppily speaking, in case of convergence,

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x_{\infty}=A^{-1} r_{0}+A^{-1} F_{\infty} z_{\infty}=A^{-1}\left(r_{0}+F_{\infty} z_{\infty}\right)
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Proving convergence is the hard task.

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\mathcal{R}_{k}(z) \equiv \frac{\chi_{k}(z)}{\chi_{k}(0)}=1-\mathcal{L}_{k}^{0}\left[1-\delta_{z 0}\right](z)=\operatorname{det}\left(I_{k}-z C_{k}^{-1}\right) .
$$

- Generalization:

$$
\mathcal{R}_{l+1: k}(z) \equiv \frac{\chi_{l+1: k}(z)}{\chi_{l+1: k}(0)}=1-\mathcal{L}_{l+1: k}^{0}\left[1-\delta_{z 0}\right](z) . \quad l=0,1, \ldots, k .
$$

Two types of polynomials $\Rightarrow$ two expressions for the QOR residuals.

## Residual polynomials and QOR residuals

## Theorem (QOR residuals)

Suppose $q_{1}=r_{0} /\left\|r_{0}\right\|$ and let all $C_{l+1: k}$ be invertible. Let $x_{k}$ denote the QOR iterate and $r_{k}=r_{0}-A x_{k}$ the corresponding residual.
Then

$$
\begin{align*}
r_{k} & =\mathcal{R}_{k}(A) r_{0}+\sum_{l=1}^{k} z_{l k} \mathcal{L}_{l+1: k}^{0}\left[1-\delta_{z 0}\right](A) f_{l}  \tag{5}\\
& =\mathcal{R}_{k}(A) r_{0}-\sum_{l=1}^{k} z_{l k} \mathcal{R}_{l+1: k}(A) f_{l}+F_{k} z_{k} .
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First expression: related to perturbation amplification.

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\end{align*}
$$

First expression: related to perturbation amplification.
Second expression: related to the attainable accuracy.

## An example: Lanczos' method

We used the diagonal matrix

$$
A=\operatorname{diag}([\operatorname{linspace}(0,1,50), 3])
$$

and the starting vector

$$
e=\operatorname{ones}(51,1)
$$

in an implementation of Lanczos' method in MATLAB on a PC conforming to ANSI/IEEE 754 with machine precision eps (1) $=2^{-52} \approx 2.2204 \cdot 10^{-16}$.

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Additionally, we heavily used the symbolic toolbox, i.e., MAPLE.

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